

A Systematic Analysis of the Juniper Dual EC Incident

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the grugq
@thegrugq

 Follow

Woah! Juniper discovers a backdoor to decrypt VPN traffic (and remote admin) has been inserted into their OS source forums.juniper.net/t5/Security-In...

2:09 PM - 17 Dec 2015

  1,653  654

- Administrative Access (CVE-2015-7755)
- VPN Decryption (CVE-2015-7756)

Administrative Access Backdoor

```
ADD      R0, R5, #0x44
LDR      R1, =aSUnSU ; "<<< %s(un='%s') = %u"
BL       strcmp
CMP      R0, #0
BNE      loc_13DC78
MOV      R0, #0xFFFFFFFF
LDMDB   R11, {R4-R8,R11,SP,PC}
```

Extra check in auth_admin_internal allows admin login using password:

```
<<< %s(un='%s') = %u
```

Changed constants in an H.D. Moore diff

P-256 Weierstraß b

5AC635D8AA3A93E7B3EBBD5576CC53B0F63BCE3C3E27D2604B
6B17D1F2E12C4247F8BCE6E563A440F277037D812DEB33A0F
FFFFFFFF00000000FFFFFFFFFFFFFFFFBCE6FAADA7179E84F3B9CAC2FC632551

P-256 P x coord

P-256 field order

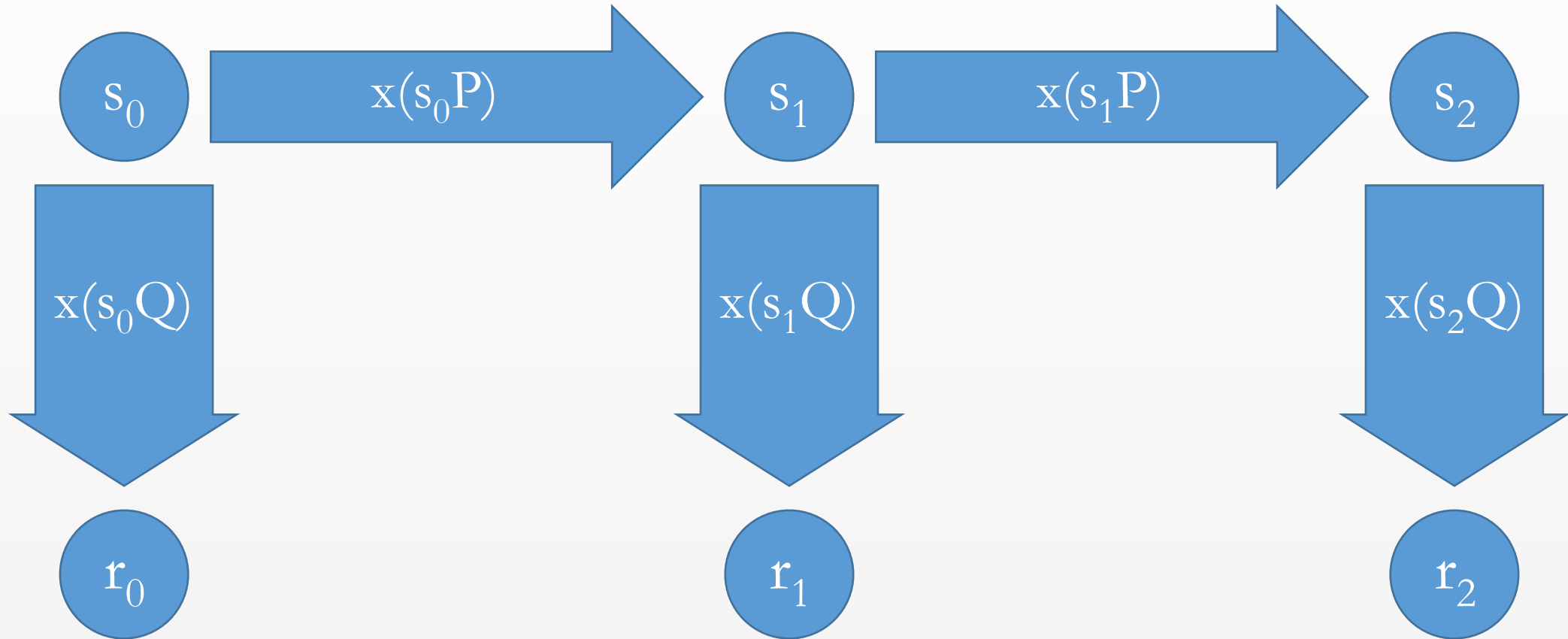
bad: 9585320EEAF81044F20D55030A035B11BECE81C785E6C933E4A8A131F6578107
good: 2c55e5e45edf713dc43475effe8813a60326a64d9ba3d2e39cb639b0f3b0ad10
nist: c97445f45cdef9f0d3e05e1e585fc297235b82b5be8ff3efca67c59852018192

Reverse engineering shows changed values are x coords for Dual EC point Q

Dual EC DRBG History

- Early 2000s: Created by the NSA and pushed towards standardization
- 2004: Published as part of ANSI x9.82 part 3 draft
- 2004: RSA makes Dual EC the default CSPRNG in BSAFE (\$10mil)
- 2005: Standardized in NIST SP 800-90 draft
- 2007: Shumow and Ferguson demonstrate theoretical backdoor attack
- 2013: Snowden documents lead to renewed interest in Dual EC
- 2014: Practical attacks on TLS using Dual EC demonstrated
- 2014: NIST removes Dual EC from list of approved PRNGs

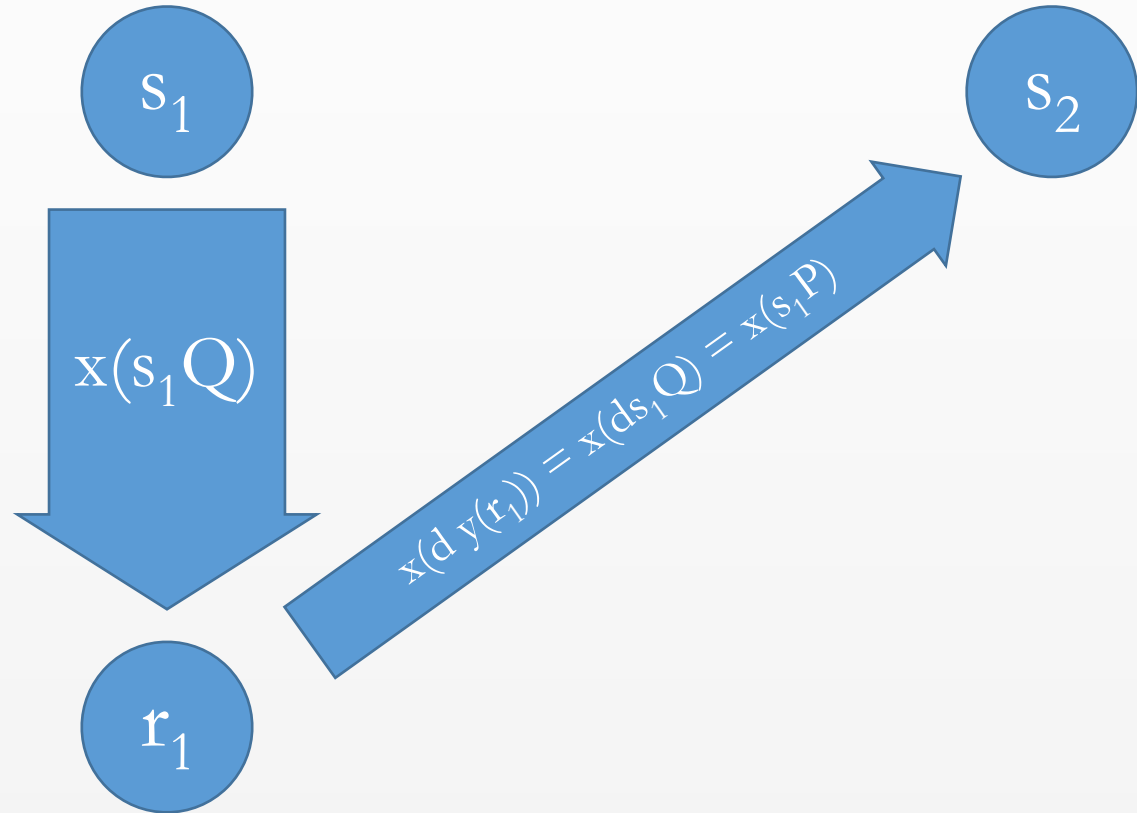
Dual EC DRBG



Note: r is actually the 30 least significant bytes of the x value

Dual EC DRBG Backdoor

Assume an attacker who
knows $\log_Q P$ aka
 d st. $P = dQ$



How to find $\log_Q P$

Disclaimer: Without more information, given P and Q there is not a way to tell if they were generated safely

- Solve the discrete log problem
- Be in charge of the official curve parameters
 - Fix Q , d , define $P = dQ$
 - Fix P , e , define $Q = eP$, compute $d = e^{-1}$
- Use your own curve parameters

e.g. the NSA

Juniper's use of Dual EC

- ScreenOS is only FIPS validated for ANSI x9.31, not Dual EC

The following product families do utilize Dual_EC_DRBG, but do not use the pre-defined points cited by NIST:

1. ScreenOS*

* ScreenOS does make use of the Dual_EC_DRBG standard, but is designed to not use Dual_EC_DRBG as its primary random number generator. ScreenOS uses it in a way that should not be vulnerable to the possible issue that has been brought to light. Instead of using the NIST recommended curve points it uses self-generated basis points and then takes the output as an input to FIPS/ANSI X.9.31 PRNG, which is the random number generator used in ScreenOS cryptographic operations.

Questions

- Why does a change in Q result in a passive VPN Decryption vulnerability?
- We doesn't Juniper's use of X9.31 protect their system against a compromise of Q?
- What is the history of the PRNG code in ScreenOS?
- How was Juniper's Q value generated?
- Is the version of ScreenOS with Juniper's Q vulnerable to attack?

We can explore the answers to these questions
using forensic reverse engineering


ScreenOS RNG

```
void prng_generate(void) {
    int time[2];
    time[0] = 0;
    time[1] = get_cycles();
    prng_output_index = 0;
    ++blocks_generated_since_reseed;
    if (!one_stage_rng())
        prng_reseed();
    for (; prng_output_index <= 0x1F; prng_output_index += 8) {
        // FIPS checks removed for clarity
        x9_31_generate_block(time, prng_seed, prng_key, prng_block);
        // FIPS checks removed for clarity
        memcpy(&prng_temporary[prng_output_index], prng_block, 8);
    }
}
```

Note that identifiers such as function and variable names are not present in the binary; we assigned these names based on analysis of the apparent function of each symbol

ScreenOS RNG

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void prng_generate(void) {
    int time[2];
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ScreenOS RNG

```
void prng_reseed(void) {  
    blocks_generated_since_reseed = 0;  
    if (dualec_generate(prng_temporary, 32) != 32)  
        error_handler("FIPS ERROR: PRNG failure, unable to reseed\n", 11);  
    memcpy(prng_seed, prng_temporary, 8);  
    prng_output_index = 8;  
    memcpy(prng_key, &prng_temporary[prng_output_index], 24);  
    prng_output_index = 32;  
}
```

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    prng_output_index = 32;  
}
```

Generate Dual EC Output

Copy to prng internals

ScreenOS RNG

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        memcpy(&prng_temporary[prng_output_index], prng_block, 8);  
    }  
}
```



Generate output with new key

Note that identifiers such as function and variable names are not present in the binary; we assigned these names based on analysis of the apparent function of each symbol


ScreenOS RNG

```
void prng_generate(void) {
    int time[2];
    time[0] = 0;
    time[1] = get_cycles();
    prng_output_index = 0;
    ++blocks_generated_since_reseed;
    if (!one_stage_rng())
        prng_reseed();
    for (; prng_output_index <= 0x1F; prng_output_index += 8) {
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        x9_31_generate_block(time, prng_seed, prng_key, prng_block);
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        memcpy(&prng_temporary[prng_output_index], prng_block, 8);
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```

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    for (; prng_output_index <= 0x1F; prng_output_index += 8) {  
        // FIPS checks removed for clarity  
        x9_31_generate_block(time, prng_seed, prng_key, prng_block);  
        // FIPS checks removed for clarity  
        memcpy(&prng_temporary[prng_output_index], prng_block, 8);  
    }  
}
```



The diagram consists of two orange callout boxes with white text. The first box, labeled 'Global Variable', has a pointer pointing to the line 'prng_output_index = 0;'. The second box, labeled 'Always true', has a pointer pointing to the line 'if (!one_stage_rng())'.

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    prng_output_index = 32;  
}
```

Global Variable

Set to 32

ScreenOS RNG

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        prng_reseed();  
    for (; prng_output_index <= 0x1F; prng_output_index += 8) {  
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        x9_31_generate_block(time, prng_seed, prng_key, prng_block);  
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        memcpy(&prng_temporary[prng_output_index], prng_block, 8);  
    }  
}
```



Never Runs



Uses same buffer

Note that identifiers such as function and variable names are not present in the binary; we assigned these names based on analysis of the apparent function of each symbol

Internet Key Exchange (IKE) protocol

- Used to establish keys for VPN session
- Two major versions, IKEv1 and v2
- Both use two phases:
 - Phase 1 establishes keys to encrypt the phase 2 handshake
 - Phase 2 establishes keys for IPSec (or other encapsulated protocol)
- Both phases present nonces and use a Diffie-Hellman key exchange

IKE Phase 1 Handshake

- Header
- Payload: Security Association
 - Contains details about which cipher suites to use
- Payload: Key Exchange
 - Contains DH key exchange data, g^x
- Payload: Nonce
 - Contains 8-128 byte random value
- Other payloads
 - Vendor info, identification, etc.

IKE Phase 1 Handshake

- Header
- Payload: Security Association
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 - Contains 8-128 byte random value
- Other payloads
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ScreenOS x comes directly from Dual EC



ScreenOS uses 32-byte nonce from Dual EC

ScreenOS Multiple Handshake Key Recovery Attack

Key Exchange value generated before Nonce means we need to see multiple handshakes

$$s_0 \rightarrow r_0$$

↓

$$s_1 \rightarrow r_1$$

↓

$$s_2 \rightarrow r_2$$

↓

$$s_3 \rightarrow r_3$$

ScreenOS Multiple Handshake Key Recovery Attack

Key Exchange value generated before Nonce means we need to see multiple handshakes

$$s_0 \rightarrow r_0$$

↓

$$s_1 \rightarrow r_1$$

↓

$$s_2 \rightarrow r_2$$

↓

$$s_3 \rightarrow r_3$$

}

IKE Handshake 1

$$KE = g^{r_0}$$

$$\text{Nonce} = r_1$$

}

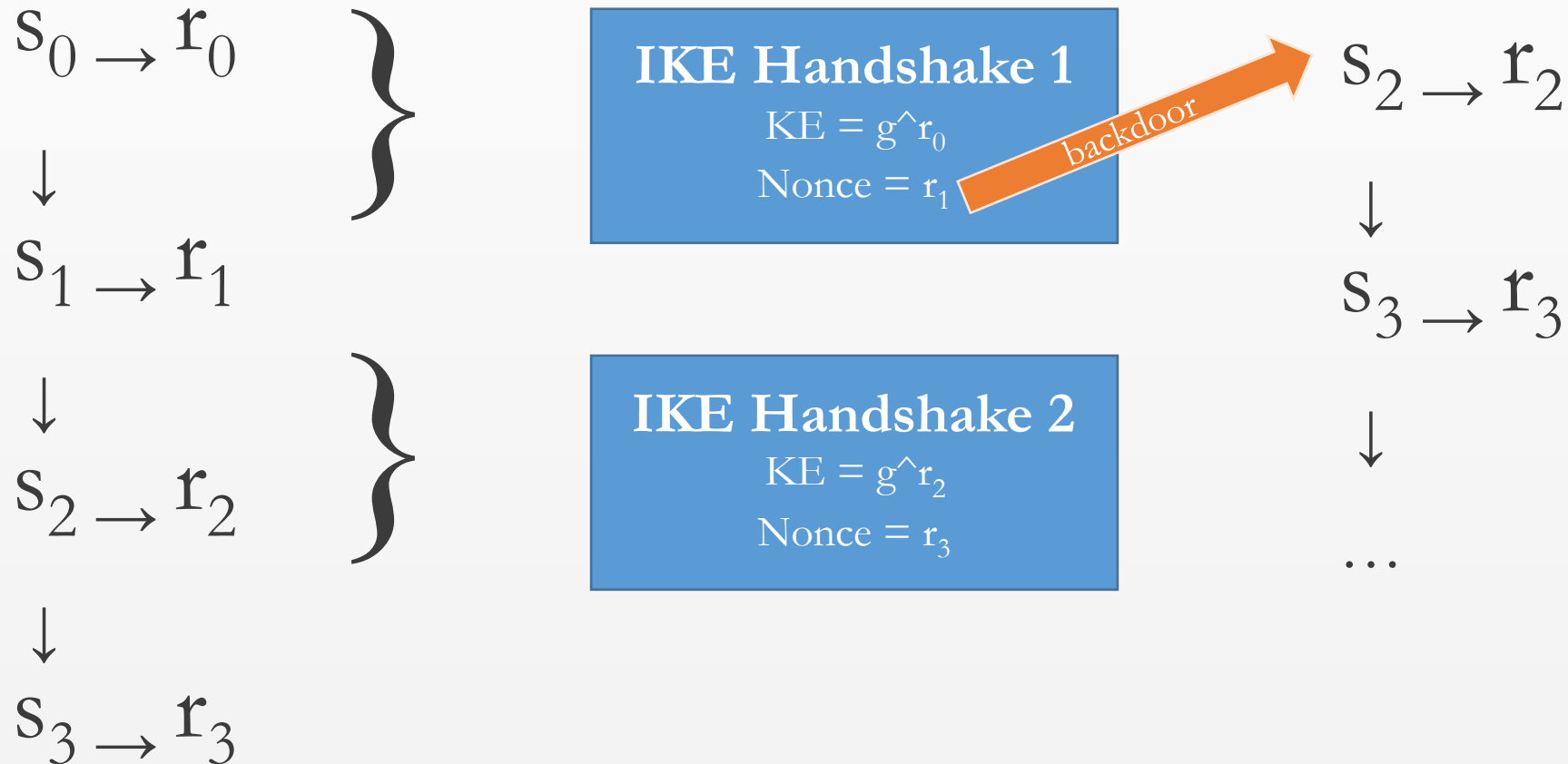
IKE Handshake 2

$$KE = g^{r_2}$$

$$\text{Nonce} = r_3$$

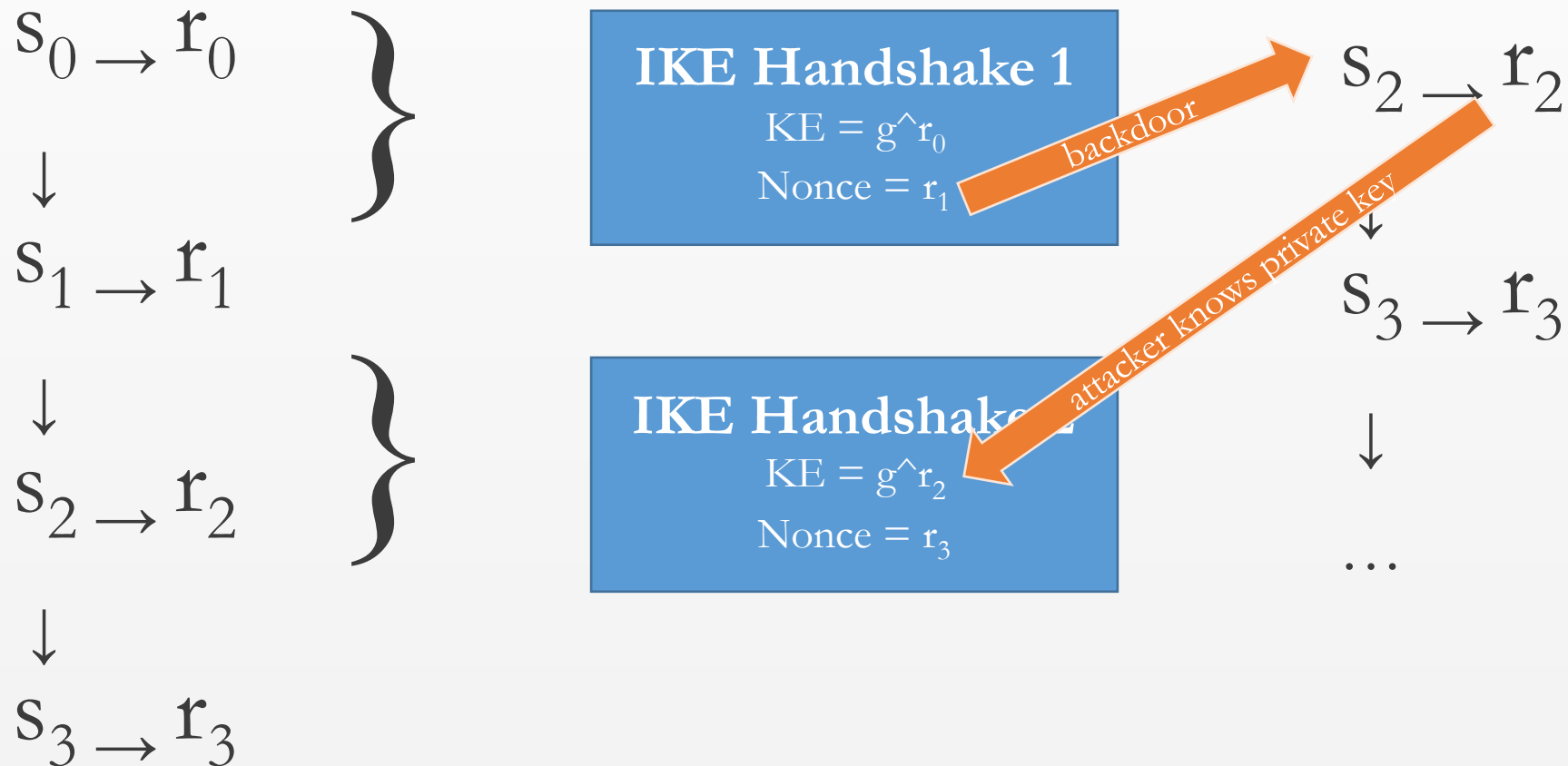
ScreenOS Multiple Handshake Key Recovery Attack

Key Exchange value generated before Nonce means we need to see multiple handshakes



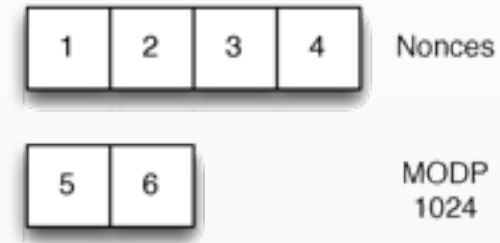
ScreenOS Multiple Handshake Key Recovery Attack

Key Exchange value generated before Nonce means we need to see multiple handshakes



Nonce Queues

- There are queues for each of:
 - Nonces
 - MODP DH groups
 - 768, 1024, 1536, and 2048 bit
 - ECP DH groups
 - 256 and 384 bit
- Filled in background process
- Nonces **always** generated before keys



At system startup



After a DH exchange

ScreenOS Single Handshake Key Recovery Attack

$s_0 \rightarrow r_0$

↓

$s_1 \rightarrow r_1$

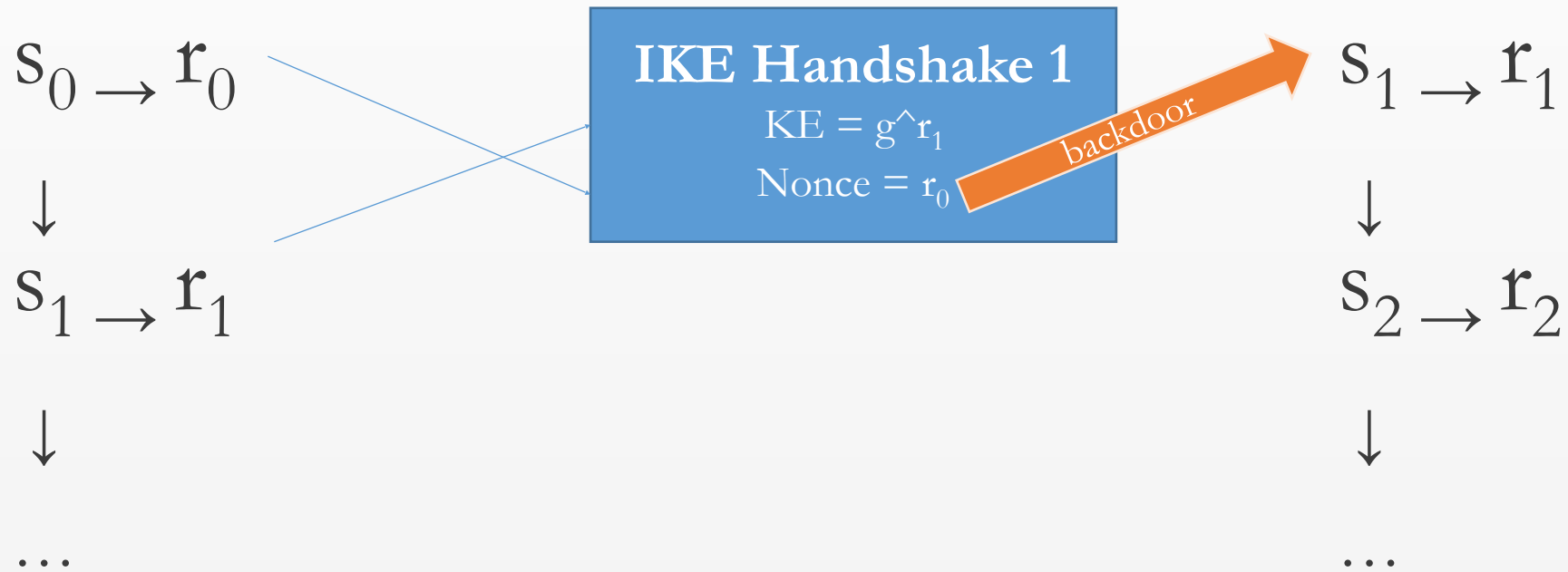
↓

...

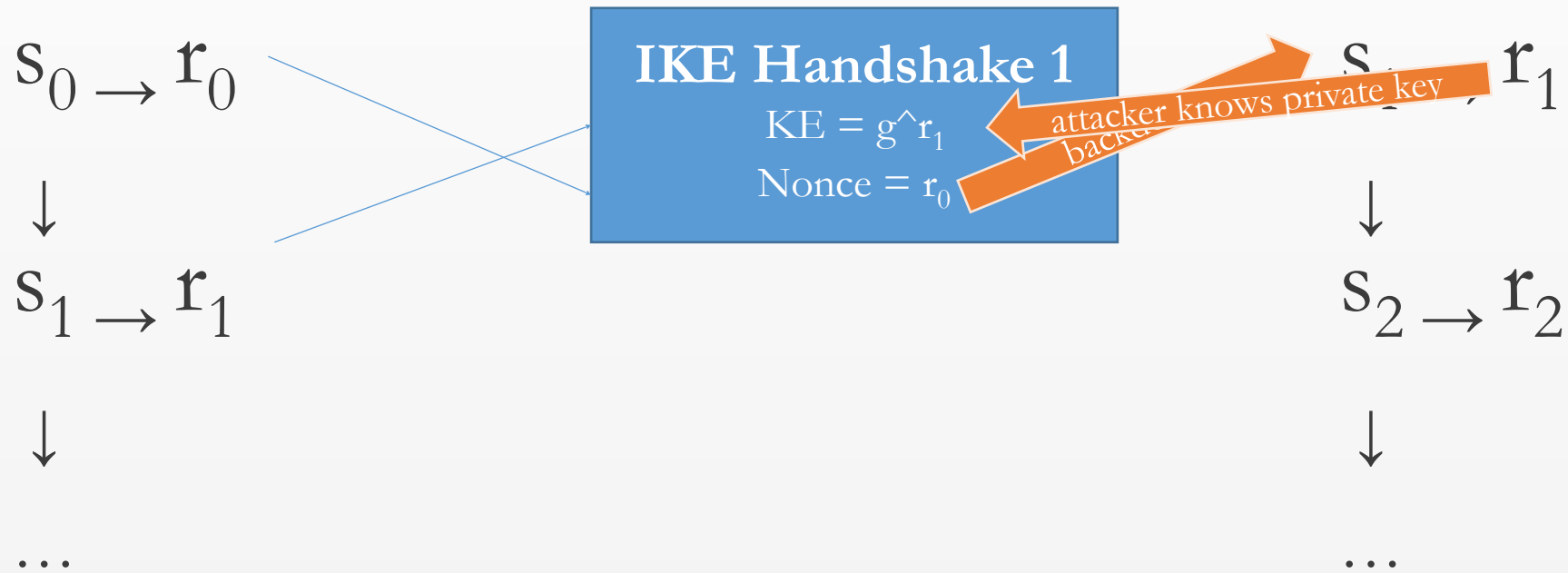
ScreenOS Single Handshake Key Recovery Attack



ScreenOS Single Handshake Key Recovery Attack



ScreenOS Single Handshake Key Recovery Attack



Caveats

- Many scenarios can downgrade single handshake attack to multiple handshake attack:
 - Fast connections exhaust queue
 - Non-DH phase 2 exchanges
 - Multiple DH queues at different rates (figure 2 in the paper)

Proof of Concept

- Purchased a Netscreen SSG 550M
- Created a modified firmware with our own Q (for which we know the discrete log d)
- Generated VPN connections in several configurations
 - IKEv1 with PSK
 - IKEv1 with RSA cert
 - IKEv2 with PSK



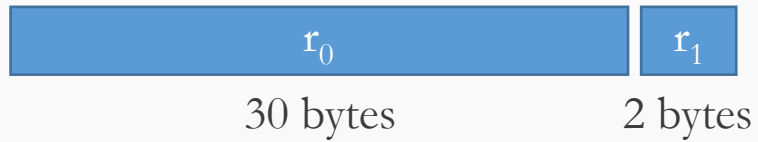
Did it Work?

- Attack worked on:
 - ~~IKEv1 with PSK~~ (attacker needs PSK)
 - IKEv1 with RSA cert
 - IKEv2 with PSK
 - Should work on IKEv2 with cert

Version History

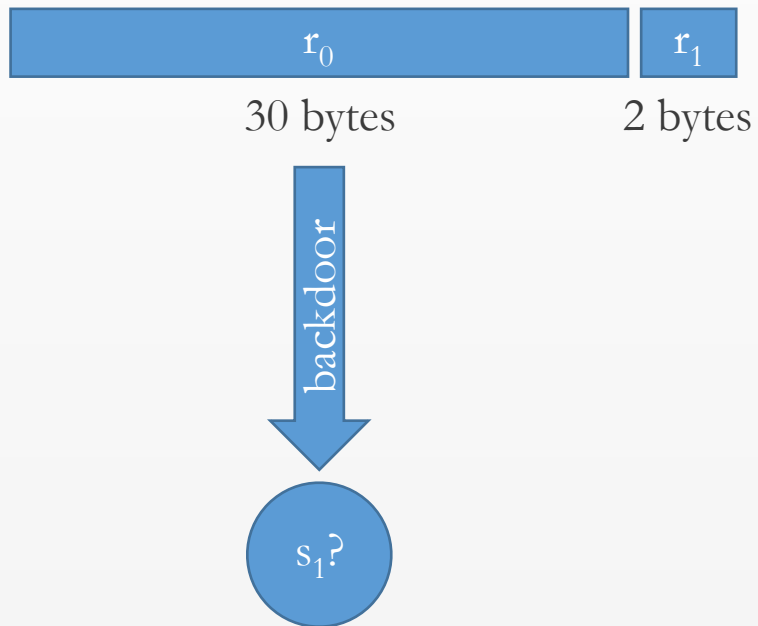
- ScreenOS 6.1.0r7 (last 6.1 version)
 - ANSI x9.31
 - Seeded by Interrupts
 - Reseeds every 10k calls
 - DH Queues
 - 20-byte IKE nonces
- ScreenOS 6.2.0r0 (first 6.2 version)
 - DualEC → ANSI x9.31
 - Reseed Bug exposes DualEC
 - Reseeds every call
 - Nonce Queues before DH Queues
 - 32-byte nonces

32-Byte Nonces



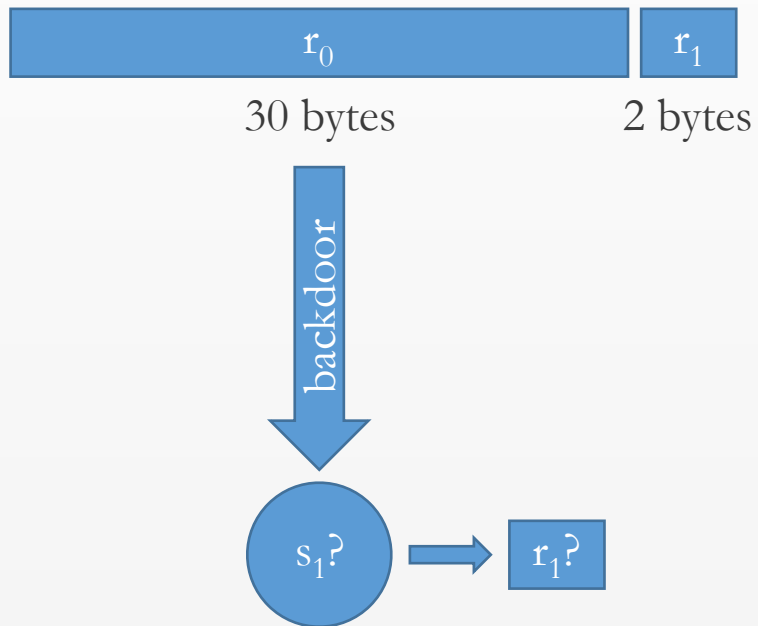
- 32-byte Dual EC outputs actually facilitate the attack:
 - Use first 30 bytes to recover 2^{15} possible states
 - For each possible state, generate a value and test against last 2 bytes

32-Byte Nonces



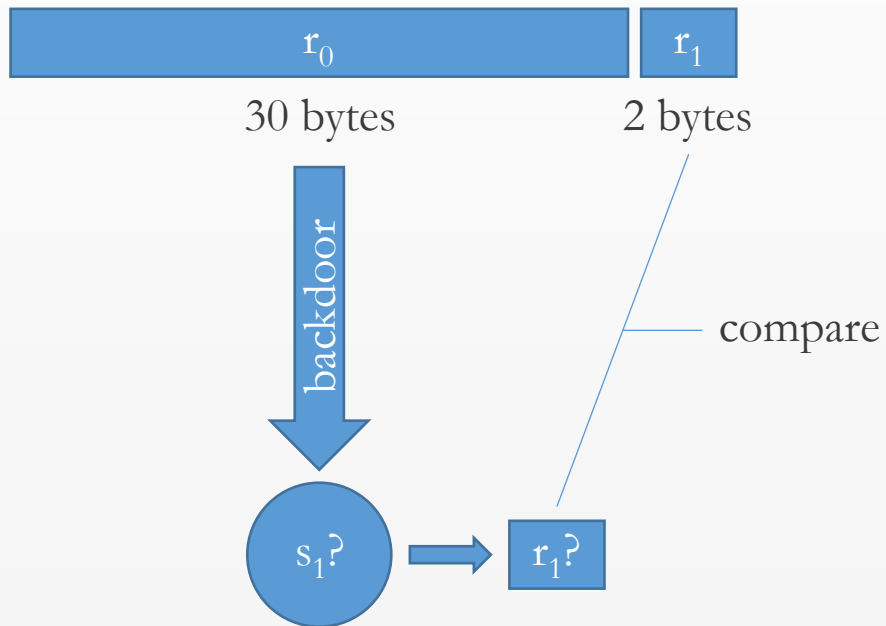
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32-Byte Nonces



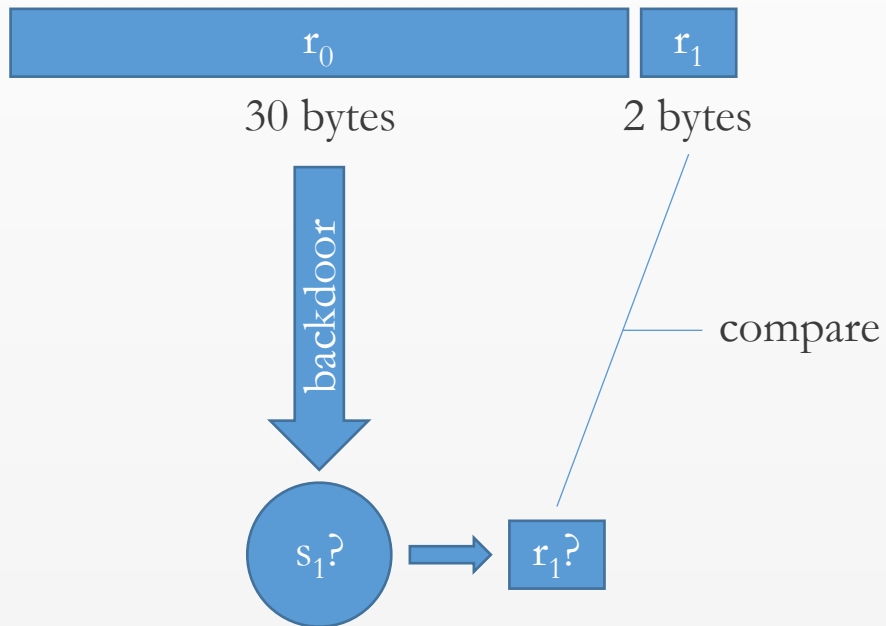
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32-Byte Nonces



- 32-byte Dual EC outputs actually facilitate the attack:
 - Use first 30 bytes to recover 2^{15} possible states
 - For each possible state, generate a value and test against last 2 bytes
- Results in 1-3 possible states in practice
- Attack is impractical with 20-byte nonce

Version History

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Attacker changed constant in 6.2.0r15

5AC635D8AA3A93E7B3EBBD55769886BC651D06B0CC53B0F63BCE3C3E27D2604B
6B17D1F2E12C4247F8BCE6E563A440F277037D812DEB33A0F4A13945D898C296
FFFFFFFF00000000FFFFFFFFFFFFFFFFBCE6FAADA7179E84F3B9CAC2FC632551

bad: 9585320EEAF81044F20D55030A035B11BECE81C785E6C933E4A8A131F6578107

good: 2c55e5e45edf713dc43475effe8813a60326a64d9ba3d2e39cb639b0f3b0ad10

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 - Nonce Queues before DH Queues
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Completely Passive Attack *Enabled in single point release*

Juniper's "fix" was to reinstate their original Q value. After our work, they removed Dual EC completely.

Answers

- ~~Why does a change in Q result in a passive VPN Decryption vulnerability?~~
- ~~We doesn't Juniper's use of X9.31 protect their system against a compromise of Q?~~
- ~~What is the history of the PRNG code in ScreenOS?~~
- How was Juniper's Q value generated?
- ~~Is the version of ScreenOS with Juniper's Q vulnerable to attack?~~

Questions?

ScreenOS Timeline

- 6.1.0r7 – ANSI generator
- 6.2.0r1 – DualEC with bugs and Juniper's Q
- 6.2.0r15 – Q changed to unknown attacker's value (12 Sept. 2012)
- 6.3.0r17 – SSH Backdoor introduced (25 April 2014?)
- 6.3.0r19b and 6.3.0.r12b – Rebuilt with backdoors removed (Dec. 2015)
- 6.3.0r22 – Dual EC removed and replaced (April 2016)

